

CSCI 3230 Data Structures

Linear Data Structures

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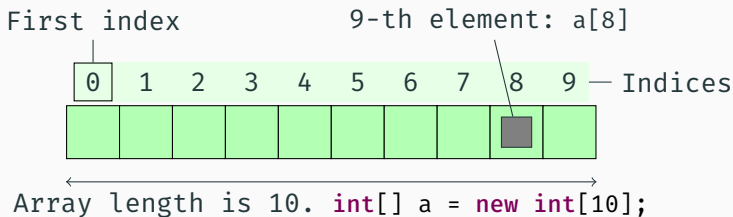
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Arrays

Array

Array: a fixed-size sequential collection of elements of the same type



- array **size is fixed** at the time of array's construction
- array elements are placed **contiguously in memory**
- indexing into a component takes **$O(1)$ time**

Array: continue ...

- **Java arrays are homogeneous**
 - all components must be of the same (object or primitive) type.
 - but, an array of an object type can contain objects of any respective subtype
- **An array is itself an object**
 - it is allocated dynamically by means of **new**
 - it is automatically deallocated when no longer referred to
- When an array is first created, **all values are automatically initialized** with
 - **0**, for an array of `int[]` or `double[]` type
 - **false**, for a `boolean[]` array
 - **null**, for an array of objects

Array: Limitations

Limitations of Arrays:

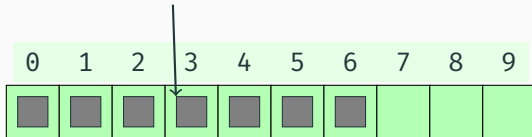
- Static data structure
- Insertion / deletion in an array is time consuming ($O(n)$, why?)

Array: Limitations

Limitations of Arrays:

- Static data structure
- Insertion / deletion in an array is time consuming ($O(n)$, why?)

insert an entry e

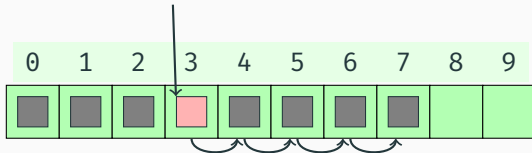


Array: Limitations

Limitations of Arrays:

- Static data structure
- Insertion / deletion in an array is time consuming ($O(n)$, why?)

insert an entry e



ArrayLists

ArrayList

ArrayList: dynamic sized arrays in Java that implement List interface.

```
1  ArrayList {  
2      size();  
3      isEmpty();  
4      add(int l, E e);  
5      get(int i);  
6      remove(int i);  
7      set(int l, E e);  
8  }
```

ArrayList

ArrayList: dynamic sized arrays in Java that implement List interface.

```
1 ArrayList {  
2     size();  
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4     add(int l, E e);  
5     get(int i);  
6     remove(int i);  
7     set(int l, E e);  
8 }
```

In the array based implementation

- The space used by the data structure is $O(n)$
- **size**, **isEmpty**, **get** and **set** run in $O(1)$ time
- **add** and **remove** run in $O(n)$ time as these two operations need to shift forward and backward on the array.

ArrayList: add(element)

When the array is full, instead of throwing an exception, we could replace the array with a larger one.

```
1  Algorithm add(element):  
2  if n = N then  
3      newArray = new array of size newN  
4      for i = 0 to N - 1 do  
5          newArray[i] = array[i]  
6      array = newArray  
7      n = n + 1  
8      array[n] = element
```

ArrayList: add(element)

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```

- **Incremental strategy:** increase the size by a constant c
- **Doubling strategy:** double the size

Incremental Strategy Analysis (optional)

n	N	
0	0	
1	c	Extend array
2	c	
\vdots	\vdots	
c	c	
c + 1	2c	Extend array
\vdots	\vdots	
2c	2c	
2c + 1	3c	Extend array
\vdots	\vdots	
n	kc	$(k = \lceil n/c \rceil)$

Replace the array $k = \lceil n/c \rceil$ times

The total time $T(n)$ is proportional to

$$n + c + 2c + 3c + 4c + \dots + kc = n + ck(k+1)/2$$

Since c is a constant,

$T(n)$ is $O(n + k^2)$, i.e., $O(n^2)$

The amortized time of an **add(element)** operation is $O(n)$

Doubling Strategy Analysis (optional)

n	N	
0	0	
1	1	Extend array
2	2	Extend array
3	4	Extend array
4	4	
5	8	Extend array
6	8	
7	8	
\vdots	\vdots	
n	2^k	$(k = \lceil \log n \rceil)$

Replace the array $k = \lceil \log n \rceil$ times

The total time $T(n)$ is proportional to

$$n + 1 + 2 + 4 + 8 + \dots + 2^k = n + 2^{k+1} - 1 \leq 5n$$

Thus $T(n)$ is $O(n)$

The amortized time of an **add(element)** operation is $O(1)$

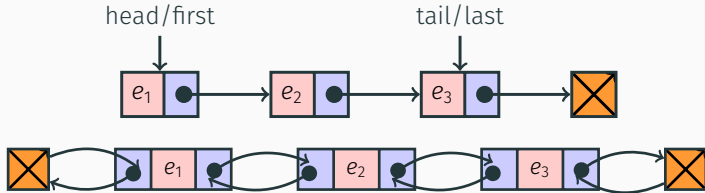
Linked Lists

Linked Lists

A concrete data structure consisting of a sequence of nodes

Each **node** stores

- element
- link to the next/previous node



Singly Linked List vs Doubly Linked List

[SinglyLinkedList.java](#), [DoublyLinkedList.java](#)

Linked Lists: Node class

```
1  private static class Node<E> {
2      private E element;
3      private Node<E> next;
4      // private Node<E> prev;
5      public Node(E e, Node<E> n) {
6          element = e;
7          next = n;
8          // prev = p;
9      }
10     public E getElement() { return element; }
11     public Node<E> getNext() { return next; }
12     // public Node<E> getPrev() { return prev; }
13     public void setNext(Node<E> n) { next = n; }
14     // public void setPrev(Node<E> p) { prev = p; }
15 }
```

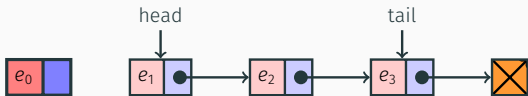
Linked Lists

Singly Linked List

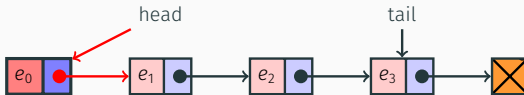
Singly Linked List: ADT

```
1  private Node<E> head = null;
2  private Node<E> tail = null;
3  /** Number of nodes in the list */
4  private int size = 0;
5
6  public int size()
7  public boolean isEmpty()
8  public E first()
9  public E last()
10 public void addFirst(E e)
11 public void addLast(E e)
12 public E removeFirst()
```

Singly Linked List: `addFirst(E e)`, `removeFirst()`

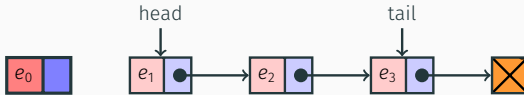


Singly Linked List: addFirst(E e), removeFirst()



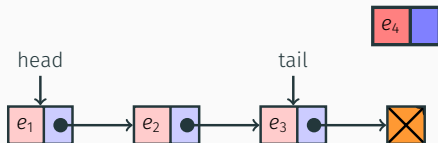
```
1 public void addFirst(E e) {  
2     head = new Node<>(e, head);  
3     if (size == 0)  
4         tail = head;  
5     size++;  
6 }
```

Singly Linked List: addFirst(E e), removeFirst()

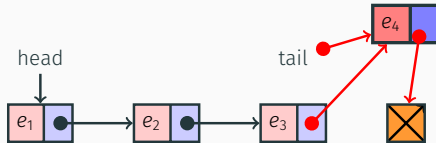


```
1 public E removeFirst() {  
2     if (isEmpty()) return null;  
3     E answer = head.getElement();  
4     head = head.getNext();  
5     size --;  
6     if (size == 0)  
7         tail = null;  
8     return answer;  
9 }
```

Singly Linked List: addLast(E e), removeLast()

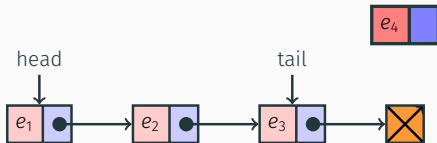


Singly Linked List: addLast(E e), removeLast()



```
1 public void addLast(E e) {  
2     Node<E> newest = new Node<>(e, null);  
3     if (isEmpty()) // previously empty list  
4         head = newest;  
5     else // new node after existing tail  
6         tail.setNext(newest);  
7     tail = newest; // new node becomes the tail  
8     size++;  
9 }
```

Singly Linked List: addLast(E e), removeLast()



There is **no constant-time way** to update the tail to point to the previous node

Why?

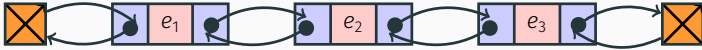
Linked Lists

Doubly Linked List

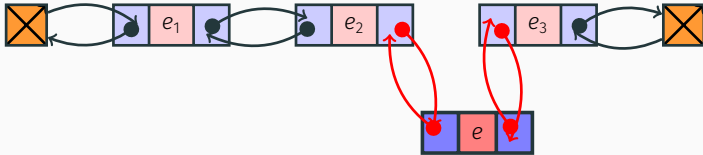
Doubly Linked List: ADT

```
1  private Node<E> head = null;
2  private Node<E> tail = null;
3  private int size = 0;
4
5  public int size()
6  public boolean isEmpty()
7  public E first()
8  public E last()
9  public void addFirst(E e)
10 public void addLast(E e)
11 public E removeFirst()
12 public E removeLast()
13 private void addBetween
14     (E e, Node<E> predecessor, Node<E> successor)
15 private E remove(Node<E> node)
```

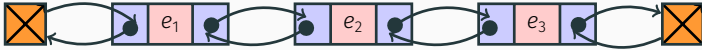
Doubly Linked List: addBetween and remove



Doubly Linked List: addBetween and remove

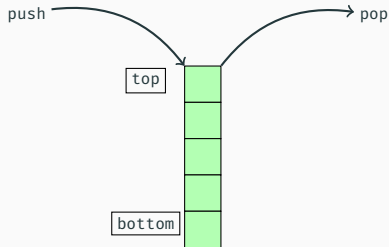


Doubly Linked List: addBetween and remove



Stacks

Stack



- The Stack ADT stores **arbitrary objects**
- Insertions and deletions follow the **last-in first-out** scheme

Stack Interface

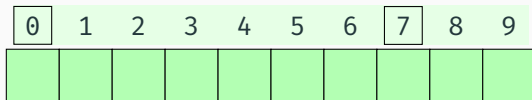
```
1  public interface Stack<E> {  
2      int size();  
3      boolean isEmpty();  
4      void push(E e);  
5      E top();  
6      E pop();  
7  }
```

Note: `java.util.Stack` provides push and pop, but differs in other respects.

[ArrayStack.java](#), [LinkedStack.java](#)

Array-based Stack: ArrayStack.java

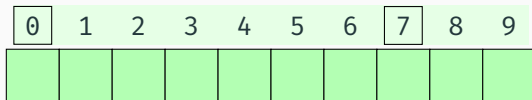
```
1  public class ArrayStack<E> implements Stack<E> {  
2      private E[] data;  
3      private int t = -1;  
4      // index of the top element in stack  
5      ...  
6      public int size()  
7      public boolean isEmpty()  
8      public void push(E e)  
9      public E pop()  
10 }
```



Which side is “bottom”?

Array-based Stack: ArrayStack.java

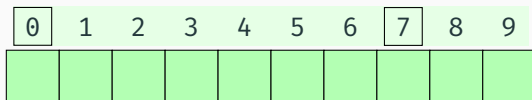
```
1  public class ArrayStack<E> implements Stack<E> {  
2      private E[] data;  
3      private int t = -1;  
4      // index of the top element in stack  
5      ...  
6      public int size()  
7      public boolean isEmpty()  
8      public void push(E e)  
9      public E pop()  
10 }
```



Which side is “bottom”? bottom top

Array-based Stack: ArrayStack.java

```
1  public class ArrayStack<E> implements Stack<E> {  
2      private E[] data;  
3      private int t = -1;  
4      // index of the top element in stack  
5      ...  
6      public int size()  
7      public boolean isEmpty()  
8      public void push(E e)  
9      public E pop()  
10 }
```

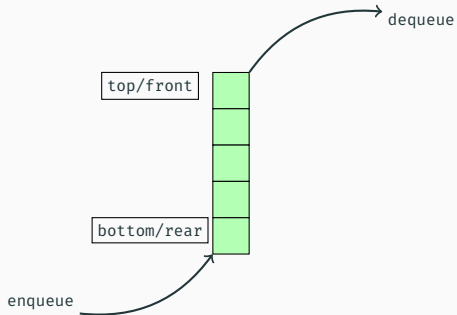


Which side is “bottom”? **bottom** **top**

- **Performance:** each operation runs in time $O(1)$
- **Limitations:** fixed size

Queues

Queue



- The Queue ADT stores **arbitrary objects**
- Insertions and deletions follow the **first-in first-out** scheme

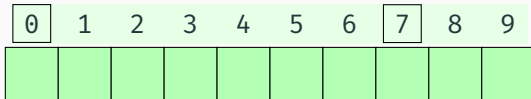
Queue Interface

```
1 public interface Queue<E> {  
2     int size();  
3     boolean isEmpty();  
4     void enqueue(E e);  
5     E first();  
6     E dequeue();  
7 }
```

ArrayQueue.java, LinkedList.java

Array-based Queue

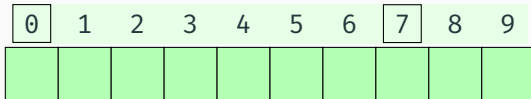
```
1  public class ArrayQueue<E> implements Queue<E> {  
2      /** Generic array used for storage of queue elements. */  
3      private E[] data;  
4      /** Index of the top element of the queue in the array. */  
5      private int f = 0;  
6      /** Current number of elements in the queue. */  
7      private int sz = 0;  
8      ...  
9      public int size()  
10     public boolean isEmpty()  
11     public void enqueue(E e) throws IllegalStateException  
12     public E first()  
13     public E dequeue()  
14 }
```



Which side is “front”?

Array-based Queue

```
1  public class ArrayQueue<E> implements Queue<E> {  
2      /** Generic array used for storage of queue elements. */  
3      private E[] data;  
4      /** Index of the top element of the queue in the array. */  
5      private int f = 0;  
6      /** Current number of elements in the queue. */  
7      private int sz = 0;  
8      ...  
9      public int size()  
10     public boolean isEmpty()  
11     public void enqueue(E e) throws IllegalStateException  
12     public E first()  
13     public E dequeue()  
14 }
```



Which side is “front”? front rear

Array-based Queue



mod operation

```
1 Algorithm enqueue(o):  
2   if size() = N - 1 then  
3     throw  
      IllegalStateException  
4   else  
5     r = (f + sz) mod N  
6     Q[r] = o  
7     sz = (sz + 1)
```

```
1 Algorithm dequeue():  
2   if isEmpty() then  
3     return null  
4   else  
5     o = Q[f]  
6     f = (f + 1) mod N  
7     sz = (sz - 1)  
8     return o
```

- **Performance:** each operation runs in time $O(1)$
- **Limitations:** fixed size

Thank you!

Questions?